Managing Fragmented Database in Distributed Database Environment

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Abstract: One of the mechanisms for managing data is replication since it improves data access and reliability. However, the amount of various data grows rapidly since technology is widely available at a low-cost. Problem arises when the transactions only work with a part of the database, not as a whole. This paper will discuss about a technique called database fragmentation. Fragmentation in distributed database is very useful in terms of usage, reliability and efficiency. Handling fragmented database replication becomes a challenging issue to administrators since the distributed database is scattered into split replica partitions or fragments. We address how to build a reliable system by using a distributed database fragmentation algorithm. The result shows that managing database replication with fragmentation will help a transaction to work with the specific part of the database rather than the entire database.

Keywords: Distributed Systems, Distributed Database, Centralized Database, Database Fragmentation

1 Introduction

Organizations critically need to supply recent data to users who may be geographically remote, while at the same time handle a volume request of data. Data replication is a way to deal with this problem since it provides users with fast, local access to shared data and protects availability of applications because alternate data access options exist (1, 2, 3). Distributed database replication involves the process of copying and maintaining database objects in multiple databases that make up a Distributed Database Systems (DDS) (4, 5).

In case of database failures, the total system of centralized databases comes to a halt. However, in distributed systems, when a component fails, the functioning of the system continues may be at a reduced performance. Hence distributed systems are more reliable. However, a distributed Database Management Systems (DBMS) that hides the distributed nature from the user, works more complex than a centralized DBMS. The fact that data can be replicated also adds an extra level of complexity to the distributed DBMS. If the software does not handle data replication adequately, there will be degradation in availability, reliability and performance compared to the centralized system. In addition, increased in complexity means the procurement and maintenance costs for a distributed DBMS also to be higher than centralized DBMS. Hence, the storage and communication cost are some of the important issues to be addressed in order to allow distributed users to efficiently and safely access data from many different sites.

2 Overview Database Fragmentation

Distributed database is a logically interrelated collection of shared data physically distributed over a computer network. Database can be decomposed into numbers of fragments. Fragmentation plays an important role in distributed database systems because usually, applications work with only some of...
the relations rather than entire of it (6, 7). In data distribution, it is better to work with subsets of relations as the unit of distribution. The other benefit from fragmentation is efficiency. Data is stored close to where it is most frequently used and for data that is not needed, it is not stored. By using fragmentation, a transaction can be divided into several subqueries that operate on fragments. So, it will increase the degrees of parallelism. Besides, it is also good for security as data not required for local applications is not stored. So, it will not be available to unauthorized users. There are mainly three type of fragmentation called horizontal, vertical and mixed.

A Horizontal Fragmentation

Horizontal fragmentation groups together the tuples in a relation that are used by the important transactions (7, 9). A horizontal fragment is produced by specifying a predicate that performs a restriction on the tuples in the relation. It is defined using the Selection operation of relational algebra. Given a relation $R$, a horizontal fragment is defined as $\sigma_p (R)$, where $p$ is a predicate based on one or more attributes of the relation.

<table>
<thead>
<tr>
<th>CustomerID</th>
<th>CustomerName</th>
<th>CustomerCity</th>
<th>Gender</th>
</tr>
</thead>
<tbody>
<tr>
<td>201998</td>
<td>Ain</td>
<td>Kota Bharu</td>
<td>Female</td>
</tr>
<tr>
<td>201923</td>
<td>Auni</td>
<td>Kuala Terengganu</td>
<td>Female</td>
</tr>
<tr>
<td>202039</td>
<td>Flo</td>
<td>Kota Bharu</td>
<td>Male</td>
</tr>
<tr>
<td>202102</td>
<td>Arya</td>
<td>Kuantan</td>
<td>Male</td>
</tr>
</tbody>
</table>

Figure 1 shows the illustration of horizontal fragmentation. Horizontal fragmentation involves forming a subset of the global relation by selecting tuples based on the value of one or more attributes which are also called scan attributes. The type of horizontal fragmentation is known as Primary if the scan attributes are contained in the relation being fragmented. However if the scan attributes are contained in a relation which owns the relation being fragmented, the type of horizontal fragmentation is called Secondary. The benefits of horizontal fragmentation are extensive when fragments can be formed such that query or update transactions for a relation at a site is largely localized to the fragment at that site (7, 9). During horizontal fragmentation, rules to ensure completeness, reconstruction and disjointness must be followed (7, 9, 10, 11). Horizontal fragmentation is most effective if the locality of reference can be established by using site-based attributes for fragmentation. It improves response time, availability and concurrency of transactions.

B Vertical Fragmentation

Vertical fragmentation groups together the attributes in a relation that are used jointly by the important transactions (7). A vertical fragment is defined using the Projection operation of the
Managing Fragmented Database in Distributed Database Environment

relational algebra. Given a relation $R$, a vertical fragmentation is defined as $\Pi a_1, \ldots, a_n (R)$, where $a_1, \ldots, a_n$ are attributes of the relation $R$, where $n$ is equal to the number of the attributes.

### Figure 2: Vertical Fragmentation

![Table 1](image1.png)

- **CustomerID** | **CustomerName** | **CustomerCity** | **Gender**
- 201998 | Ain | Kota Bharu | Female
- 201923 | Auni | Kuala Terengganu | Female
- 202039 | Flo | Kota Bharu | Male
- 202102 | Arya | Kuantan | Male

![Table 2](image2.png)

- **CustomerID** | **CustomerName** | **Gender**
- 201998 | Ain | Female
- 201923 | Auni | Female
- 202039 | Flo | Male
- 202102 | Arya | Male

![Table 3](image3.png)

- **CustomerID** | **CustomerCity**
- 201998 | Kota Bharu
- 201923 | Kuala Terengganu
- 202039 | Kota Bharu
- 202102 | Kuantan

Figure 2 shows the illustration of vertical fragmentation. Vertical fragmentation involves dividing the attributes of a relation into groups and then projecting the relation over each group. The attributes should be no overlapping across the groups with the exception of primary key attributes or tuple ID substantial when fragments can be formed such that query or update transactions for a relation at a site are largely localized to the fragment at that site is the benefit of vertical fragmentation. Vertical fragmentation comes generally with side-effects and requires a detailed analysis for establishing its benefits such as improving response time, availability and concurrency of transactions. The additional overhead of implicit joins will have to be incurred for queries spanning multiple fragments. The use of the implicit joins is practical when different applications located at multiple sites share data and to merge the two partitioning approaches to result in mixed partitioning (7, 9). All the updates are executed on the master copy while other copies which are called snapshots are refreshed occasionally. Queries are routed to a local copy, if available. This practice improves availability and concurrency of query operations and it does not degrade response time for update operations. Vertical fragmentation requires consistency analysis at the application level and it is most suitable for networks with high delays and low reliability.

### C Hybrid Fragmentation

For some applications horizontal or vertical fragmentation of a database schema by itself is insufficient to adequately distribute the data. Instead, mixed or hybrid fragmentation is required.
Figure 3 shows the illustration of hybrid fragmentation. A hybrid fragment is defined using the Selection and Projection operations of the relational algebra. Given a relation $R$, a mixed fragment is defined as $\sigma_p (\Pi a_1, \ldots, a_n(R))$ or $\Pi a_1, \ldots, a_n(\sigma_p (R))$ where $p$ is a predicate based on one or more attributes of $R$ and $a_1, \ldots, a_n$ are attributes of $R$.

### 3 Database Fragmentation in Distributed Database Systems

In this section, we proposed database fragmentation in distributed database systems. The following notations are defined:

i. $S$ is relation in database.
ii. $x$ is an instant which will be modified
iii. $y$ is an instant which will not be modified
iv. $S_1$ is a vertical fragmented relation with instant $x$.
v. $S_2$ is a vertical fragmented relation without instant $x$.
vi. $P_k$ is a primary key.
vii. $P_k,x$ is a primary key with data $x$.
viii. $P_k,y$ is a primary with data $y$, where $y \neq x$
ix. $S_1(P_k,x)$ and $S_1(P_k,y)$ are a horizontal fragmentation relation

In distributed database systems, each primary replica copies other database to its replicas. The fragmentation process starts with relation $S$ fragmented into $S_1$ and $S_2$ using vertical fragmentation. Next, the fragmented relation, $S_1$ will be fragmented again but this time using horizontal fragmentation into $S_1(P_k,x)$ and $S_1(P_k,y)$ leaving only instant for primary key and data that need to be updated.

### 4 Distributed Database Fragmentation Example Case Study

Let’s say, Transaction 1 ($T_1$), requests to update data Debit for user with UserID: GS001 from RM1500 to RM3000.
Relation S fragmented into $S_1$ and $S_2$ using vertical fragmentation as shown in Table 3 and Table 4.

$S_1 = \text{UserID (primary key), Debit}$
$S_2 = \text{UserID (primary key), Credit, Date}$

Relation $S_1$ then fragmented using horizontal fragmentation into $S_{1(\text{Pk},x)}$ and $S_{1(\text{Pk},y)}$. After the fragmentation, $T_i$ will update data $x$ and only will replicate $S_{1(\text{Pk},x)}$ to the other replicas.
5 Conclusion

Handling fragmented database replication is very important in order to preserve the data reliability with low communication cost. Therefore, a database fragmentation in distributed database systems has been proposed to manage the fragmented database replication. From the example case, it shows that the fragmentation is able to reduce the size of the table to smaller subsets and make them available with less network access time. Furthermore, fragmentation may help to reduce storage space because this data will be distributed evenly among replica servers instead of saving all data to all servers.

References
